## Department of Mathematics

Faculty of Science
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## Discrete Mathematics

## Yarmouk University

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## Section 1.1

## Propositions.

DEFINITION:
A proposition is a declarative sentence to which we can assign a truth value of either true or false but not both.
Examples:

1. 4 is even number (true) (1).
2. 5 is even number (false) (0).

We will use A, B, C, P, Q, R, to present propositions there letters called 1ogical variables.

- Compound proposition.

We will use connectives to form compound propositions.

1. To connectives AND denoted by $\wedge$.
2. To connectives OR denated by $\underline{V}$.
3. To connectives NOT denoted by $\sim$.

- $\quad \mathrm{A}, \mathrm{B}$ variables:

1. $(A \wedge B) \ldots$ Conjunction of $A$ and $B$.
2. $(A \vee B) \ldots$ Disjunction of $A$ and $B$.

Examples:

1. Read this statements and answer following questions.

| questions | 1 (true) |
| :--- | :--- |
| A: the earth is round. | 0 (false) |
| B: the sun is cold. | 1 (true) |
| C: it rains in Spain. | Answers |
| questions 0 (false) <br> $\mathrm{A} \wedge \mathrm{B}:$ the earth is round and the sun is cold. 1 (true) <br> $\mathrm{B} \vee \mathrm{C}$ : the surn is cold or it rains in Spain. 1 (true) <br> $\sim \mathrm{B}$ : the sun is not cold.  |  |

P and $\sim \mathrm{P}$ always have apposite truth value.
Truth table ( $2^{1}$ ).

| $\mathbf{P}$ | $\sim \mathbf{P}$ |
| :---: | :---: |
| 1 | 0 |
| 0 | 1 |

- $\quad A$ and $B$ and $A \wedge B$.

Truth table ( $2^{2}$ ).

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{A} \wedge \mathbf{B}$ |
| :---: | :---: | :---: |
| 1 | 1 | 1 |
| 1 | 0 | 0 |
| 0 | 1 | 0 |
| 0 | 0 | 0 |

- $\quad A$ and $B$ and $A \vee B$.

Truth table ( $2^{2}$ ).


- $\quad \mathrm{A} \wedge(\mathrm{B} \vee \mathrm{C})$.

Truth table ( $2^{3}$ ).

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{B} \vee \mathbf{C}$ | $\mathbf{A} \wedge(\mathbf{B} \vee \mathbf{C})$ |
| :---: | :---: | :---: | :---: | :---: |
| 1 | $\mathbf{1}$ | 1 | 1 | 1 |
|  |  | 1 | 0 | 1 |
|  | 0 | 1 | 1 | 1 |
|  |  | 0 | 0 | 0 |

- If two forms have identical truth table then said to be equivalent.
- If A and B are equivalent forms, we write $\mathrm{A}=\mathrm{B}$ and called logically identical.
- LAWS:
- DeMorgan's laws.

1. $\sim(A \wedge B)=\sim A \vee \sim B$
2. $\sim(A \vee B)=\sim A \wedge \sim B$

## PROOF:

1. $\sim(\mathrm{A} \wedge \mathrm{B})=\sim \mathrm{A} \vee \sim \mathrm{B}$

Truth table ( $2^{2}$ ).

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{A} \wedge \mathbf{B}$ | $\sim(\mathbf{A} \wedge \mathbf{B})$ | $\sim \mathbf{A}$ | $\sim \mathbf{B}$ | $\sim \mathbf{A} \vee \sim \mathbf{B}$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |
| 1 | 0 | 0 | $\mathbf{1}$ | 0 | 1 | $\mathbf{1}$ |
| 0 | 1 | 0 | $\mathbf{1}$ | 1 | 0 | $\mathbf{1}$ |
| 0 | 0 | 0 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| $\sim(\mathrm{~A} \vee \mathrm{~B})$ | $=\sim \mathrm{A} \wedge \sim \mathrm{B}$ |  |  |  |  |  |

2. $\sim(\mathrm{A} \vee \mathrm{B})=\sim \mathrm{A} \wedge \sim \mathrm{B}$

Truth table ( $2^{2}$ ).


- Distributive laws

1. $A \wedge(B \vee C)=(A \wedge B) \vee(A \wedge C)$
2. $A \vee B \wedge C)=(A \vee B) \wedge(A \vee C)$

## PROOF:

1. $A \wedge(B \vee C \perp(A \wedge B) \vee(A \wedge C)$

Truth table ( $2^{3}$ )

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{B} \vee \mathbf{C}$ | $\mathbf{\wedge}(\mathbf{B} \vee \mathbf{C})$ | $\mathbf{A} \wedge \mathbf{B}$ | $\mathbf{A} \wedge \mathbf{C}$ | $(\mathbf{A} \wedge \mathbf{B}) \vee(\mathbf{A} \wedge \mathbf{C})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 |  | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 1 | 1 | 0 | 1 | $\mathbf{1}$ | 1 | 0 | $\mathbf{1}$ |
| 1 | 0 | 1 | 1 | $\mathbf{1}$ | 0 | 1 | $\mathbf{1}$ |
| 1 | 0 | $\mathbf{0}$ | 0 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |
| $\mathbf{Q}$ | 1 | 1 | 1 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |
| 0 | 1 | 0 | 1 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |
| 0 | 0 | 1 | 1 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |
| 0 | 0 | 0 | 0 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |

2. $\quad \mathrm{A} \vee(\mathrm{B} \wedge \mathrm{C})=(\mathrm{A} \vee \mathrm{B}) \wedge(\mathrm{A} \vee \mathrm{C})$

Truth table $\left(2^{3}\right)$

| $\mathbf{A}$ | $\mathbf{B}$ | $\mathbf{C}$ | $\mathbf{B} \wedge \mathbf{C}$ | $\mathbf{A} \vee(\mathbf{B} \wedge \mathbf{C})$ | $\mathbf{A} \vee \mathbf{B}$ | $\mathbf{A} \vee \mathbf{C}$ | $(\mathbf{A} \vee \mathbf{B}) \wedge(\mathbf{A} \vee \mathbf{C})$ |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 1 | 1 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 1 | 1 | 0 | 0 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 1 | 0 | 1 | 0 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 1 | 0 | 0 | 0 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 0 | 1 | 1 | 1 | $\mathbf{1}$ | 1 | 1 | $\mathbf{1}$ |
| 0 | 1 | 0 | 0 | $\mathbf{0}$ | 1 | 0 | $\mathbf{0}$ |
| 0 | 0 | 1 | 0 | $\mathbf{0}$ | 0 | 1 | $\mathbf{0}$ |
| 0 | 0 | 0 | 0 | $\mathbf{0}$ | 0 | 0 | $\mathbf{0}$ |



